The structure of graphs excluding Gem and \hat{K}_4 as induced minors

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- $\leqslant_{\mathrm{im}}:$ induced minor relation
 - vertex deletion;
 - edge contraction;
 - but no edge deletion.

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between minors and induced subgraphs

Excluding a graph



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Every H-minor-free graph can be obtained by clique sums of order at most c_H from graphs that can be c_H -nearly embedded in some surface, in which H cannot be embedded.

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 - induced matchings (excluding $k \cdot K_2$);
 - induced packings of cycles (excluding $k \cdot K_3$);
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Our contribution: structure of graphs excluding Δ / \checkmark as induced minors.







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- excluding each of them yields a well-quasi-order (wrt. \leq_{im});
- we use the decomposition theorems to prove the wqo result.

Decomposition of Gem-induced minor-free graphs

Theorem (Błasiok, Kamiński, R., Trunck, 2015+)

If Gem $\leq_{im} G$ and G is 2-c, removing some ≤ 6 vertices gives a disjoint union of cographs and (straight) paths.



(Reminder: Gem =
$$\heartsuit$$
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Decomposition of \hat{K}_4 -induced minor-free graphs

Theorem (Błasiok, Kamiński, R., Trunck, 2015+)

If $\hat{K}_4 \not\leq_{im} G$ and G is 2-c, then

- either $K_4 \not\leq_{im} G$;
- or G is a subdivision of a graph on at most 9 vertices;
- or V(G) = M ∪ C where G[M] is complete multipartite, G[C] is a cycle and adjacencies are "binary".



(Reminder:
$$\hat{K}_4 = \Delta$$
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• we assume that $\hat{K}_4 \not\leq_{im} G$ and G has a proper K_4 -subdivision S;

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- $\bullet \ \rightarrow \ {\rm cycle-multipartite} \ {\rm decomposition}$

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- if G has a K_2 -cutset, every cc is a cograph or a straight path;
- if G has a $\overline{K_2}$ -cutset: case analysis again!
- eventually, the removal of \leqslant 6 vertices gives cographs and straight paths.

2-c Gem-induced minor-free graph:



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Note: structural results on these classes of graphs have also been obtained recently by Belmonte et al..

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Thank you!