Clique-width of Restricted Graph Classes

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Aussois

Motivation

Most natural problems in algorithmic graph theory are NP-complete.

Want to find restricted classes of graphs where we can solve some problems in polynomial time.

Best if we can find classes where lots of problems can be solved in polynomial time.

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Why Clique-width?

Theorem (Courcelle, Makowsky and Rotics 2000, Kobler and Rotics 2003, Rao 2007, Oum 2008, Grohe and Schweitzer 2015)

Any problem expressible in "monadic second-order logic with quantification over vertices" (and certain other classes of problems) can be solved in polynomial time on graphs of bounded clique-width.

This includes:

- **Vertex Colouring**
- Maximum Independent Set
- Minimum Dominating Set
- Hamilton Path/Cycle
- Partitioning into Perfect Graphs
- Graph Isomorphism
- ...

Clique-width

The clique-width is the minimum number of labels needed to construct G by using the following four operations:

(i) creating a new graph consisting of a single vertex v with label i (represented by $i(v)$)

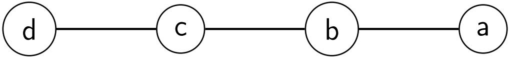
(ii) taking the disjoint union of two labelled graphs G_1 and G_2 (represented by $G_1 \oplus G_2$)

(iii) joining each vertex with label i to each vertex with label j $(i \neq j)$ (represented by $\eta_{i,j}$)

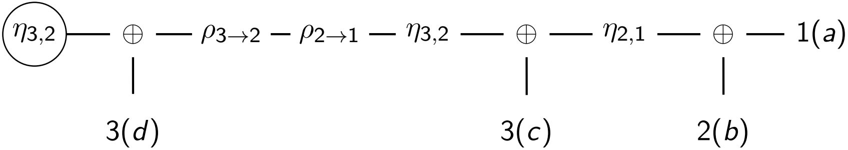
(iv) renaming label i to j (represented by $\rho_{i \to j}$)

For example, P_4 has clique-width 3.

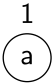
Clique-width



$$\eta_{3,2}(3(d) \oplus \rho_{3\to 2}(\rho_{2\to 1}(\eta_{3,2}(3(c) \oplus \eta_{2,1}(2(b) \oplus 1(a))))))$$



Clique-width



$1(a)$

$1(a)$

Clique-width

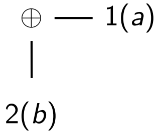
$2(b)$ $1(a)$

$1(a)$

$2(b)$

Clique-width

$2(b) \oplus 1(a)$



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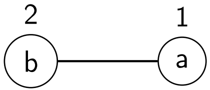
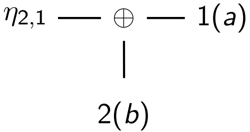


$$\eta_{2,1}(2(b) \oplus 1(a))$$



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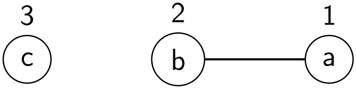
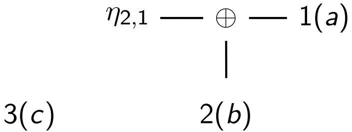


$3(c)$ $\eta_{2,1}(2(b) \oplus 1(a))$



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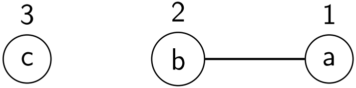
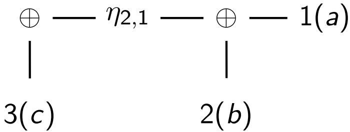


$$3(c) \oplus \eta_{2,1}(2(b) \oplus 1(a))$$



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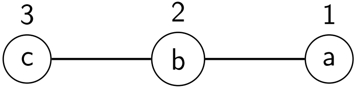
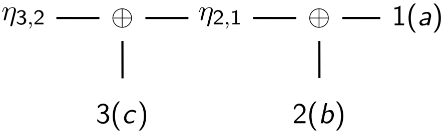


$$\eta_{3,2}(3(c) \oplus \eta_{2,1}(2(b) \oplus 1(a)))$$



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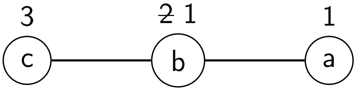
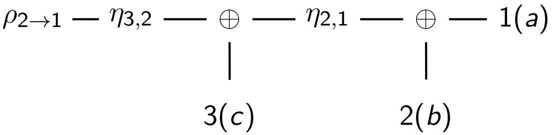


$$\rho_{2\to 1}(\eta_{3,2}(3(c) \oplus \eta_{2,1}(2(b) \oplus 1(a))))$$



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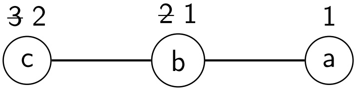
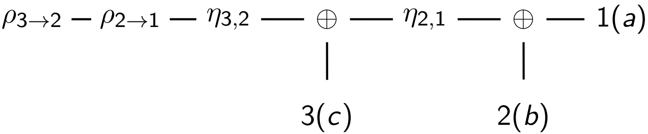


$$\color{red}{\rho_{3\to 2}}(\rho_{2\to 1}(\eta_{3,2}(3(c) \oplus \eta_{2,1}(2(b) \oplus 1(a)))))$$



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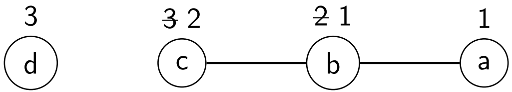
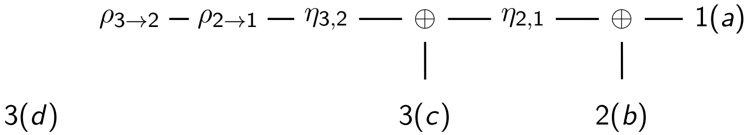


$3(d)$ $\quad \rho_{3\to 2}(\rho_{2\to 1}(\eta_{3,2}(3(c) \oplus \eta_{2,1}(2(b) \oplus 1(a)))))$



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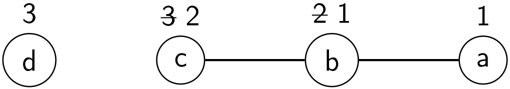

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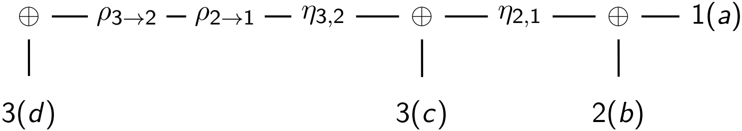

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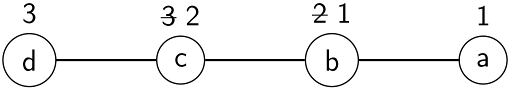

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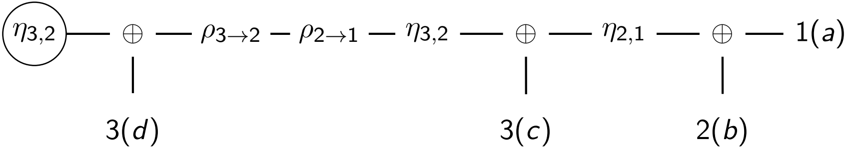

Calculating clique-width

Theorem (Fellows, Rosamond, Rotics, Szeider 2009)

Calculating clique-width is NP-hard.

Theorem (Corneil, Habib, Lanlignel, Reed, Rotics 2012)

Can detect graphs of clique-width at most 3 *in polynomial time.*

It's not known if this is also the case for graphs of clique-width 4.

Theorem (Oum 2008)

Can find a c*-expression for a graph* G *where* $c \leq 8^{\mathrm{cw}(G)} - 1$ *in cubic time.*

The clique-width of all graphs up to 10 vertices has been calculated (Heule & Szeider 2013).

Why clique-width?

- "Equivalent" to rank-width and NLC-width
- Generalises tree-width
- "Equivalent" to tree-width on graphs of bounded degree

The following operations don't change the clique-width by "too much"

- Complementation
- Bipartite complementation
- Vertex deletion
- Edge subdivision (for graphs of bounded-degree)

Need only look at graphs that are

- prime
- 2-connected

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Aim

Underlying Research Question

What kinds of graph properties ensure bounded clique-width?

Hereditary Classes

A graph H is an induced subgraph of G if H can be obtained by deleting vertices of G, written $H \subseteq_i G$.

P_4 $\quad$ $3P_1$ $\quad$ $P_1 + P_2$

So $P_1 + P_2 \subseteq_i P_4$, but $3P_1 \not\subseteq_i P_4$.

A class of graphs is hereditary if it is closed under taking induced subgraphs.

Let S be a set of graphs. The class of S-free graphs is the set of graphs that do not contain any graph in S as an induced subgraph.

For example: bipartite graphs are the $(C_3, C_5, C_7, \ldots)$-free graphs.

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H-free graphs

Theorem (D., Paulusma 2015)

The class of H-free graphs has bounded clique-width if and only if $H \subseteq_i P_4$.

Colouring H-free graphs

Theorem (Král', Kratochvíl, Tuza & Woeginger, 2001)

The Vertex Colouring *problem is polynomial-time solvable for H-free graphs if and only if $H \subseteq_i P_1 + P_3$ or P_4, otherwise it is NP-complete.*

$P_1 + P_3$

P_4

Colouring (H_1, H_2)-free graphs

The Vertex Colouring problem is polynomial-time solvable for (H_1, H_2)-free graphs if

1. H_1 or H_2 is an induced subgraph of $P_1 + P_3$ or of P_4
2. $H_1 \subseteq_i K_{1,3}$, and $H_2 \subseteq_i C_3^{++}$, $H_2 \subseteq_i C_3^*$ or $H_2 \subseteq_i P_5$
3. $H_1 \neq K_{1,5}$ is a forest on at most six vertices or $H_1 = K_{1,3} + 3P_1$, and $H_2 \subseteq_i \overline{P_1 + P_3}$
4. $H_1 \subseteq_i sP_2$ or $H_1 \subseteq_i sP_1 + P_5$ for $s \geq 1$, and $H_2 = K_t$ for $t \geq 4$
5. $H_1 \subseteq_i sP_2$ or $H_1 \subseteq_i sP_1 + P_5$ for $s \geq 1$, and $H_2 \subseteq_i \overline{P_1 + P_3}$
6. $H_1 \subseteq_i P_1 + P_4$ or $H_1 \subseteq_i P_5$, and $H_2 \subseteq_i \overline{P_1 + P_4}$
7. $H_1 \subseteq_i P_1 + P_4$ or $H_1 \subseteq_i P_5$, and $H_2 \subseteq_i \overline{P_5}$
8. $H_1 \subseteq_i 2P_1 + P_2$, and $H_2 \subseteq_i \overline{2P_1 + P_3}$ or $H_2 \subseteq_i \overline{P_2 + P_3}$
9. $H_1 \subseteq_i \overline{2P_1 + P_2}$, and $H_2 \subseteq_i 2P_1 + P_3$ or $H_2 \subseteq_i P_2 + P_3$
10. $H_1 \subseteq_i sP_1 + P_2$ for $s \geq 0$ or $H_1 = P_5$, and $H_2 \subseteq_i \overline{tP_1 + P_2}$ for $t \geq 0$
11. $H_1 \subseteq_i 4P_1$ and $H_2 \subseteq_i \overline{2P_1 + P_3}$
12. $H_1 \subseteq_i P_5$, and $H_2 \subseteq_i C_4$ or $H_2 \subseteq_i \overline{2P_1 + P_3}$.

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The Vertex Colouring problem is polynomial-time solvable for (H_1, H_2)-free graphs if

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4. $H_1 \subseteq_i sP_2$ or $H_1 \subseteq_i sP_1 + P_5$ for $s \geq 1$, and $H_2 = K_t$ for $t \geq 4$
5. $H_1 \subseteq_i sP_2$ or $H_1 \subseteq_i sP_1 + P_5$ for $s \geq 1$, and $H_2 \subseteq_i \overline{P_1 + P_3}$
6. $H_1 \subseteq_i P_1 + P_4$ or $H_1 \subseteq_i P_5$, and $H_2 \subseteq_i \overline{P_1 + P_4}$
7. $H_1 \subseteq_i P_1 + P_4$ or $H_1 \subseteq_i P_5$, and $H_2 \subseteq_i \overline{P_5}$
8. $H_1 \subseteq_i 2P_1 + P_2$, and $H_2 \subseteq_i \overline{2P_1 + P_3}$ or $H_2 \subseteq_i \overline{P_2 + P_3}$
9. $H_1 \subseteq_i \overline{2P_1 + P_2}$, and $H_2 \subseteq_i 2P_1 + P_3$ or $H_2 \subseteq_i P_2 + P_3$
10. $H_1 \subseteq_i sP_1 + P_2$ for $s \geq 0$ or $H_1 = P_5$, and $H_2 \subseteq_i \overline{tP_1 + P_2}$ for $t \geq 0$
11. $H_1 \subseteq_i 4P_1$ and $H_2 \subseteq_i \overline{2P_1 + P_3}$
12. $H_1 \subseteq_i P_5$, and $H_2 \subseteq_i C_4$ or $H_2 \subseteq_i \overline{2P_1 + P_3}$.

The class of (H_1, H_2)-free graphs has bounded clique-width if:

1. H_1 or $H_2 \subseteq_i P_4$;
2. $H_1 = sP_1$ and $H_2 = K_t$ for some s, t;
3. $H_1 \subseteq_i P_1 + P_3$ and $\overline{H_2} \subseteq_i K_{1,3} + 3P_1$, $K_{1,3} + P_2$, $P_1 + S_{1,1,2}$, P_6 or $S_{1,1,3}$;
4. $H_1 \subseteq_i 2P_1 + P_2$ and $\overline{H_2} \subseteq_i 2P_1 + P_3$, $3P_1 + P_2$ or $P_2 + P_3$;
5. $H_1 \subseteq_i P_1 + P_4$ and $\overline{H_2} \subseteq_i P_1 + P_4$ or P_5;
6. $H_1 \subseteq_i 4P_1$ and $\overline{H_2} \subseteq_i 2P_1 + P_3$;
7. $H_1, \overline{H_2} \subseteq_i K_{1,3}$.

and it has unbounded clique-width if:

1. $H_1 \notin \mathcal{S}$ and $H_2 \notin \mathcal{S}$;
2. $\overline{H_1} \notin \mathcal{S}$ and $\overline{H_2} \notin \mathcal{S}$;
3. $H_1 \supseteq_i K_{1,3}$ or $2P_2$ and $\overline{H_2} \supseteq_i 4P_1$ or $2P_2$;
4. $H_1 \supseteq_i 2P_1 + P_2$ and $\overline{H_2} \supseteq_i K_{1,3}$, $5P_1$, $P_2 + P_4$ or P_6;
5. $H_1 \supseteq_i 3P_1$ and $\overline{H_2} \supseteq_i 2P_1 + 2P_2$, $2P_1 + P_4$, $4P_1 + P_2$, $3P_2$ or $2P_3$;
6. $H_1 \supseteq_i 4P_1$ and $\overline{H_2} \supseteq_i P_1 + P_4$ or $3P_1 + P_2$.

Theorem (D., Paulusma 2015)

This leaves 13 cases where it is unknown if the clique-width of (H_1, H_2)-free graphs is bounded or not (up to some equivalence relation).

1. $H_1 = 3P_1, \overline{H_2} \in \{P_1 + P_2 + P_3, P_1 + 2P_2, P_1 + P_5, P_1 + S_{1,1,3}, P_2 + P_4, S_{1,2,2}, S_{1,2,3}\}$;
2. $H_1 = 2P_1 + P_2, \overline{H_2} \in \{P_1 + P_2 + P_3, P_1 + 2P_2, P_1 + P_5\}$;
3. $H_1 = P_1 + P_4, \overline{H_2} \in \{P_1 + 2P_2, P_2 + P_3\}$ *or*
4. $H_1 = \overline{H_2} = 2P_1 + P_3$.

Working paper with François Dross and Daniël Paulusma: The 5 cases in orange have bounded clique-width.

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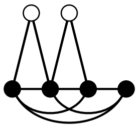
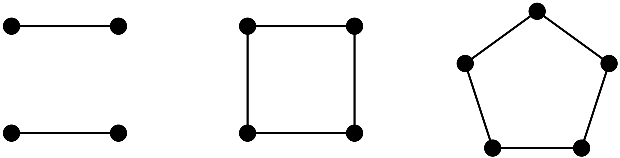
Split Graphs

A graph is split if its vertices can be partitioned into an independent set and a clique.

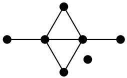
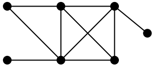


Equivalently, split graphs are the $(2P_2, C_4, C_5)$-free graphs.



$2P_2$ C_4 C_5

H-free Split Graphs

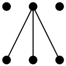
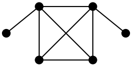
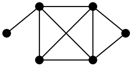
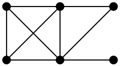
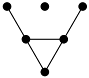
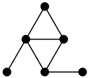
 

F_4 F_5

Theorem (Brandstädt, D., Huang, Paulusma, 2015)

Let H be a graph such that neither H nor $\overline{H}$ is in $\{F_4, F_5\}$. The class of H-free split graphs has bounded clique-width if and only if H or $\overline{H}$ is

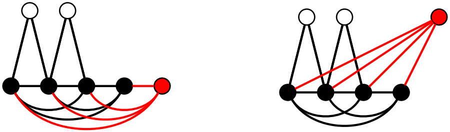
- *isomorphic to rP_1 for some $r \geq 1$ or*
- *an induced subgraph of one of:*

$K_{1,3} + 2P_1$ F_1 F_2 F_3 $\text{bull} + P_1$ Q

Split Partitions

A split graph can have more than one way of partitioning the vertices into a clique and an independent set.



Can't have two vertices both in the clique in one partition and both in the independent set in another partition.

So, if two partitions differ on two vertices, they must look like this:

These are the same! Can only differ on one vertex.

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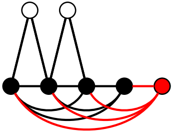
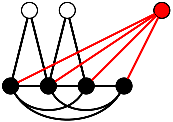

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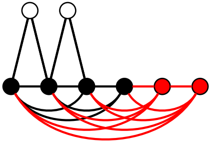
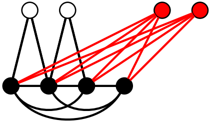

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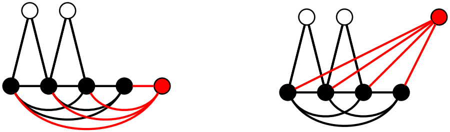

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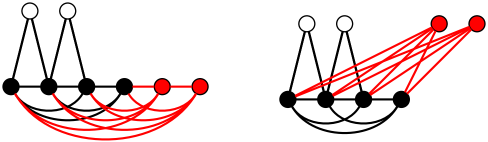

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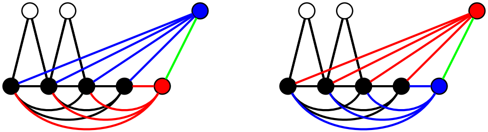

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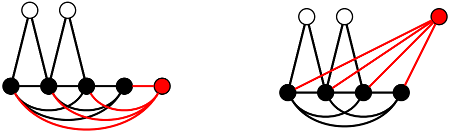

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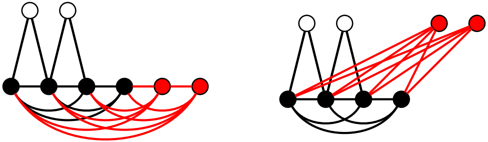

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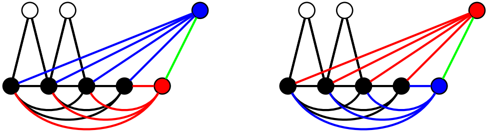

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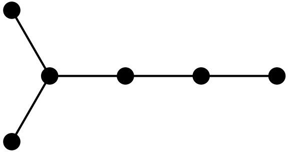
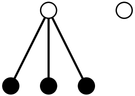
Labelled Bipartite Graphs

A labelled bipartite graph consists of a set W of white vertices and a set B of black vertices.

Let H^ℓ be a labelled bipartite graph. A bipartite graph G is weakly H^ℓ-free if it has a bipartite partition such that H is not an induced subgraph of G in a way that respects the colours.



G



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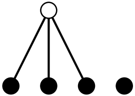

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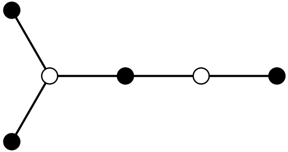

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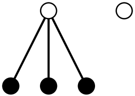

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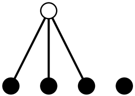

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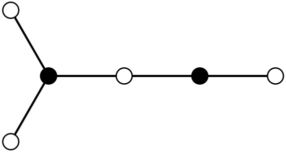

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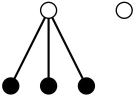

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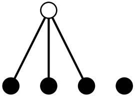
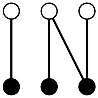


H^{ℓ^*}

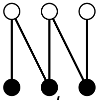
Theorem (D., Paulusma (2014))

The class of weakly H^ℓ-free bipartite graphs has bounded clique-width if and only if H^ℓ or $H^{\overline{\ell}}$ is a labelled induced subgraph of one of the following:

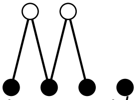
$(kP_1)^b$ for $k = 5$



$(P_2 + P_4)^b$



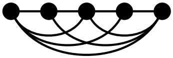
P_6^b



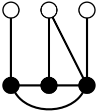
$(P_1 + P_5)^b$

Corollary

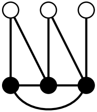
The class of H-free split graphs has bounded clique-width if H is one of the following graphs:



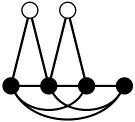
K_r for $R = 5$



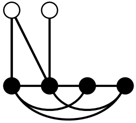
Q

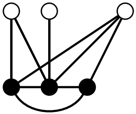


$\overline{Q}$



$\overline{\text{bull} + P_1}$





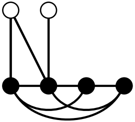
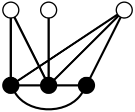
Theorem

The class of F_3-free split graphs has bounded clique-width.

Proof.

- Let G be an F_3-free split graph. Fix a split partition of G. We may assume it is prime, so no two vertices have the same neighbourhood.
- If G contains less than 19 copies of D, we can delete them and get a Q-free split graph, so done.

- Assume G contains at least 19 vertex-disjoint copies of D, say $D_1, \ldots, D_{19}$.

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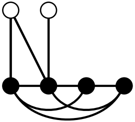 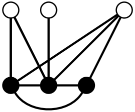

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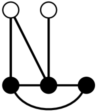

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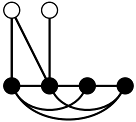 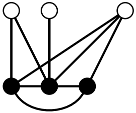

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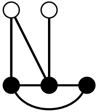

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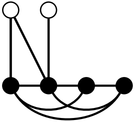

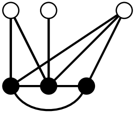

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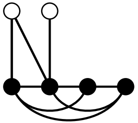

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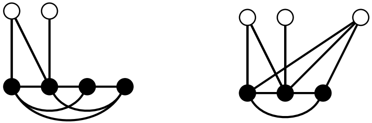

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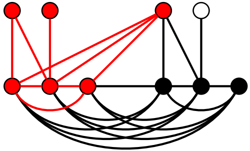

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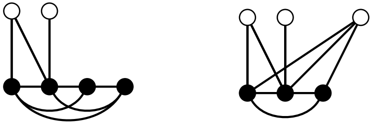

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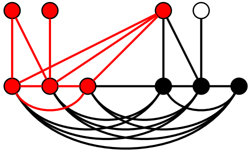

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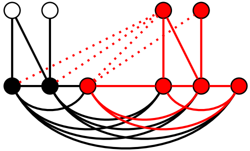

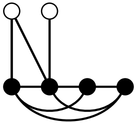

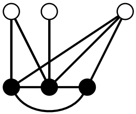

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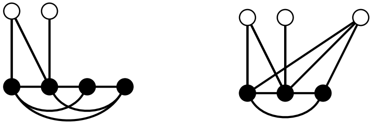
- Contradiction! □



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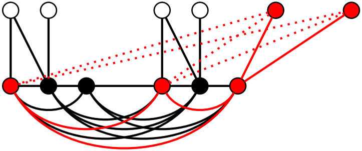

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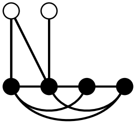

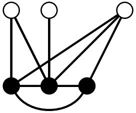

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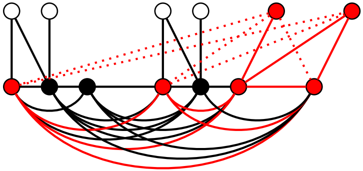

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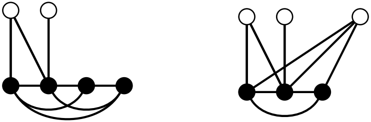

- Each white vertex in a D_i must have one of the $6 \times 6 = 36$ possible neighbourhoods in D_1 and D_2.
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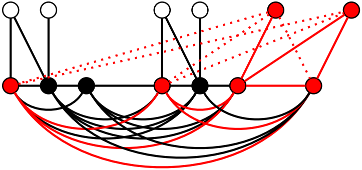
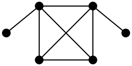
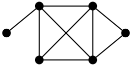


- Since no two vertices have the same neighbourhood, there must be a black vertex that distinguishes them.
- Contradiction! $\square$

H-free Chordal Graphs

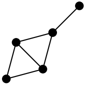


F_1

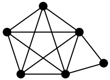


F_2

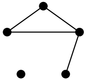
Theorem (Brandstädt, D., Huang, Paulusma 2015)
Let H be a graph with $H \notin \{F_1, F_2\}$. The class of H-free chordal graphs has bounded clique-width if and only if H is a an induced subgraph of one of:



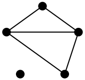
$\overline{S_{1,1,2}}$



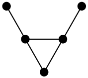
$\overline{K_{1,3} + 2P_1}$



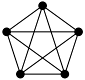
$P_1 + \overline{P_1 + P_3}$



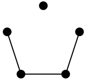
$P_1 + \overline{2P_1 + P_2}$



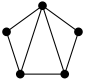
bull



K_r for $r = 5$



$P_1 + P_4$



$\overline{P_1 + P_4}$

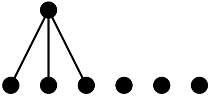
H-free Weakly Chordal Graphs

Theorem (Brandstädt, D., Huang, Paulusma 2015)
Let H be a graph. Then the class of H-free weakly chordal graphs has bounded clique-width if and only if $H \subseteq_i P_4$.

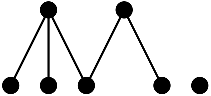
H-free Bipartite Graphs

Theorem (D., Paulusma 2014)

The class of H-free bipartite graphs has bounded clique-width if and only if H is an induced subgraph one of:

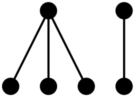


$K_{1,3} + 3P_1$

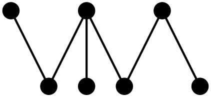


$P_1 + S_{1,1,3}$

sP_1 for some s
($s = 5$ shown)



$K_{1,3} + P_2$



$S_{1,2,3}$

Other Containment Relations

Theorem (D., Paulusma 2015)

Let $\{H_1, \ldots, H_p\}$ be a finite set of graphs. Then the following statements hold:

(i) *The class of $(H_1, \ldots, H_p)$-subgraph-free graphs has bounded clique-width if and only if $H_i \in \mathcal{S}$ for some $1 \leq i \leq p$.*

(ii) *The class of $(H_1, \ldots, H_p)$-minor-free graphs has bounded clique-width if and only if H_i is planar for some $1 \leq i \leq p$.*

(iii) *The class of $(H_1, \ldots, H_p)$-topological-minor-free graphs has bounded clique-width if and only if H_i is planar and has maximum degree at most 3 for some $1 \leq i \leq p$.*

Summary of Open Problems

For which pairs of graphs (H_1, H_2) does the class of (H_1, H_2)-free graphs have bounded clique-width? (13 open cases: see also "Clique-width of Graph Classes Defined by Two Forbidden Induced Subgraphs" D. & Paulusma, CIAC 2015 and arXiv:1405.7092.)

Upcoming result: $(3P_1, \overline{P_1 + 2P_2})$*-free graphs and four superclasses have bounded clique-width.*

For which graphs H does the class of H-free chordal graphs have bounded clique-width? (2 open cases: see also "Bounding the Clique-Width of H-free Chordal Graphs" Brandstädt, D., Huang, Paulusma, MFCS 2015 and arXiv:1502.06948.)

For which graphs H does the class of H-free split graphs have bounded clique-width? (2 open cases: see also "Bounding the Clique-Width of H-free Split Graphs" Brandstädt, D., Huang, Paulusma, Eurocomb 2015)

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Thank You!